Novel algorithms for distributed spectrum sensing

Prof. Vinod Sharma¹

Dept. of Electrical Communication Engg., IISc

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¹ joint work with Jithin K.S., Sahasranand K.R. and Febi Ibrahim

Outline

- Introduction CR and spectrum sensing
- Non parametric Sequential Hypothesis Testing algorithms
- Performance Comparison
- Asymptotic properties
- Robust tests
- Distributed set up
- Putting together what we learnt
- Performance analysis of distributed algorithm
- Asymptotics of distributed algorithm
- Conclusion

Introduction: Cognitive Radio



- CR: perform radio environment analysis, identify spectral holes and operate in those holes.
- Licensed (primary) and Unlicensed (secondary) users
- Spectrum Sensing
 - Identify spectral holes and quickly detect the onset of primary transmission
 - Primary techniques: Matched filter, Cyclostationary detector and Energy detector
 - Challenges: Shadowing, Fading and low SNR ($\sim -20dB$), sensing frequency and duration

Model for Cognitive Radio

Non parametric sequential hypothesis testing problem

• i.i.d. observations $X_i, i = 1, 2, \ldots$

$$\begin{array}{rcl} H_0 & : & X_i = N_i & & \sim P_0(\text{ p.d.f. } f_0) \\ H_1 & : & X_i = H_i S_i + N_i & & \sim P_1(\text{ p.d.f. } f_1) \end{array}$$

• Take N observations and use decision rule δ_N which

$$\min_{N,\delta_N} E[N|H_0] \text{ and } \min_{N,\delta_N} E[N|H_1],$$

subject to
$$P_{F\!A} \leq lpha$$
 and $P_{MD} \leq eta$

Issues in Spectrum sensing

- -20 dB SNR, shadowing, fading \Rightarrow Distributed detection required.
- Fast detection \Rightarrow Sequential detection preferred.
- Transmit power, channel gains, modulation, coding schemes etc. of PU transmissions are not known (SNR uncertainty)
 ⇒ (Generalised) Energy detection optimal.
- H_iS_i distribution not known \Rightarrow Non parametric/Semi parametric and composite hypothesis.
- Fast fading Rayleigh, Rician, Nakagami.
 Slow fading Log normal.
- Time varying electromagnetic interference ⇒ Distribution of SINR N_i may not be known and noise power time varying.
- Outliers present \Rightarrow Robust tests desired.
- Hence, Non parametric, distributed, robust, sequential tests preferred.

Existing methods - a discussion

• P_0 and P_1 fully known - SPRT optimal for discrete and continuous alphabet, For 0 $<\alpha,\beta<1$

$$N \stackrel{\Delta}{=} \inf\{n : \tilde{W}_n = \sum_{k=1}^n \log \frac{P_1(X_k)}{P_0(X_k)} \notin (\log \beta, -\log \alpha)\},\$$

$$\delta_N = H_1 \text{ if } \tilde{W}_N \ge -\log \alpha; H_0 \text{ if } \tilde{W}_N \le \log \beta$$

• P_0 is known, P_1 is not known - our algorithm (KTSLRT).

► Discrete alphabet case: L_n(X₁ⁿ), codelength of a universal lossless source code for X₁ⁿ. We replace W̃_n by

$$\widehat{W}_n = -L_n(X_1^n) - \log P_0(X_1^n) - n\frac{\lambda}{2}, \, \lambda > 0.$$

• λ chosen to make it Order 1 asymptotically optimal

Existing methods - a discussion

• Continuous alphabet case:

- Quantize X_k to $X_k^{\Delta} = [X_k/\Delta]\Delta$ (Uniform scalar quantization)
- Range of quantization according to f₀'s tail probabilities less than a small specific value.
- Use Universal lossless coding on $X_1^{\Delta}, X_2^{\Delta}, \dots, X_n^{\Delta}$
- Performs better than the asymptotically optimal Hoeffding test for finite alphabet sources, Kolmogorov-Smirnoff test and some other non parametric tests.

Our new algorithm - entropy test Discussion of the test

• P_0 is known, P_1 is not known.

• Replace
$$\widehat{W}_n$$
 by

$$W_n = W_{n-1} + [-\log P_0(X_n) - H(P_0) - \frac{\lambda}{2}], \ \lambda > 0.$$

where $H(P_0)$ is entropy (differential entropy for continuous case) of P_0 .

- Average drift under H_1 : $D(P_1||P_0) + H(P_1) H(P_0) \lambda/2$ Thus we take under H_1 , $\{P_1 : D(P_1||P_0) \ge \lambda, H(P_1) \ge H(P_0)\}$
- Average drift under H_0 : $-\lambda/2$.
- Does not require quantization for continuous case.
- Does not require a Universal source encoder.

Other non parametric tests

Rank Test

- i. Let $Y_i = X_i \frac{\mu_0 + \mu_1}{2}$, where X_i s are the observations.
- ii. Calculate R_i , the rank of Y_i in $Y_1, ..., Y_n$ when arranged in ascending order of their absolute values.
- iii. Test statistic is given by, $W = \sum_{i=1}^{n} Y_i \frac{R_i}{n+1}$.
- iv. Non parametric, for symmetric distributions.
- v. Use its sequential version.

Other non parametric tests

Sequential t test

- i. We extend the usual t test to make it a two sided test.
- ii. Test statistic is given by, $T_n = n \frac{\bar{X}_n \frac{\mu_0 + \mu_1}{2}}{s_n}$.

Random walk

- i. Test statistic is given by, $T_n = \sum_{i=1}^n (\bar{X}_n \frac{\mu_0 + \mu_1}{2}).$
- ii. The test statistic is iteratively computable.

Comparison with different tests: $P_0 \sim Bin(8, 0.2)$, $P_1 \sim Bin(8, 0.5)$



Laplace and Gaussian distributions example

 $\begin{array}{l} \textit{f}_1 \sim \textit{Laplace}(1,5) \\ \textit{f}_0 \sim \textit{Laplace}(0,1) \end{array}$

$$egin{aligned} & f_1 \sim \mathcal{N}(1,5) \ & f_0 \sim \mathcal{N}(0,1) \end{aligned}$$



• Entropy test outperforms well known tests in all these cases.

Case of no variance change



- Entropy test is not good when only mean difference tested.
- Rank test, t test and random walk perform much better.

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Spectrum Sensing

Energy detection

• Energy samples are computed as: $Y_j = \sum_{i=1}^{L} X_i^2$



- Entropy test performs best when there is difference in variances; naturally arises in energy detection.
- Energy detection (adding up $|X_i|^2$) is a special case of summing up $|X_i|^p$ to form test statistics. The traditional method (p = 2) is optimal for additive Gaussian noise.

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Asymptotic Properties of entropy test

Theorem
(1)
$$P_0(N < \infty) = 1$$
 and $P_1(N < \infty) = 1$.
(2) $\limsup_{t \to \infty} \frac{N_0(t)}{t} \le \frac{1}{\lambda(1/2+c)}$ a.s. and in L_1 (under H_0)
(3) $\limsup_{t \to \infty} \frac{N_1(t)}{t} \le \frac{2}{\lambda}$ a.s. and in L_1 (under H_1)
(4) $P_{FA} \stackrel{\Delta}{=} P_0(\widehat{W}_N \ge t_1) \le c_1 e^{-\Gamma^* t_1}$
where $\log E_0[e^{\Gamma^* Z_1}] = 0;$
 $Z_k = -\log P_0(X_k) - H(P_0) - \lambda/2.$
(5) $P_{MD} \stackrel{\Delta}{=} P_1(\widehat{W}_N \le t_0) \le c_0 e^{-\Gamma^* t_0}$
where $\log E_1[e^{\Gamma^* Z_1}] = 0;$

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Asymptotic Properties of t test

Theorem

(1)
$$P_0(N < \infty) = 1$$
 and $P_1(N < \infty) = 1$.

(2) Under H_i , $\frac{N(t)}{t} \rightarrow \frac{2\sigma}{\mu_1 - \mu_0}$ a.s., for i = 0, 1. Also, if $E_i[|X_1|] < \infty$ for i = 0, 1, then the convergence holds in L_1 .

(3) $P_{FA} \leq \alpha^s$ where s is the solution of $E_0[e^{s(X-\frac{\mu_0+\mu_1}{2})}] = 1$

(4)
$$P_{MD} \leq \beta^{s^*}$$
 where s^* is the solution of $E_0[e^{-s^*(X-\frac{\mu_0+\mu_1}{2})}] = 1$

Asymptotic Properties of random walk

Theorem

(1)
$$P_0(N < \infty) = 1$$
 and $P_1(N < \infty) = 1$.

(2) Under H_i , $\frac{N(t)}{t} \rightarrow \frac{2}{\mu_1 - \mu_0}$ a.s., for i = 0, 1. Also, if $E_i[|X_1|] < \infty$ for i = 0, 1, then the convergence holds in L_1 .

(3) $P_{FA} \leq \alpha^s$ where s is the solution of $E_0[e^{s(X-\frac{\mu_0+\mu_1}{2})}] = 1$

(4)
$$P_{MD} \leq \beta^{s^*}$$
 where s^* is the solution of $E_0[e^{-s^*(X-\frac{\mu_0+\mu_1}{2})}] = 1$

Robust Tests

- Use α -trimmed mean for t test and random walk.
- M test A robust version of *t* test insensitive to the tails of the distribution.

$$T_{n} = \frac{\sum_{i=1}^{n} \psi(X_{i} - \frac{(\mu_{0} + \mu_{1})}{2})}{\left(\sum_{i=1}^{n} \psi^{2}(X_{i} - \overline{X}_{n})\right)^{\frac{1}{2}}}$$

 $\psi:\mathcal{R}\mapsto\mathcal{R}\text{, non decreasing, continuous, odd and bounded function}$ Huber suggested,

$$\psi_0(z) = \begin{cases} k & , \text{if } z > k, \\ z & , \text{if } |z| \le k, \\ -k & , \text{if } z < -k. \end{cases}$$

• Applying M trimming to random walk gives

$$T_n = \sum_{i=1}^n \psi(X_i - \frac{(\mu_0 + \mu_1)}{2}).$$

This is iterative unlike α -trimmed random walk/t test or M test.

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Robust tests

Performance Comparison



Distributed set up

General algorithm



$$\begin{split} & W_{k,l} = -\log P_0(X_1^n) - nH(P_0) - n\frac{\lambda}{2}, \, \lambda > 0. \text{ (Entropy)} \\ & Y_{k,l} = b_1 \mathbb{I}\{W_{k,l} \ge -\log \alpha_l\} + b_0 \mathbb{I}\{W_{k,l} \le \log \beta_l\}. \\ & Y_k = \sum_{l=1}^{L} Y_{k,l} + Z_k. \\ & F_k = F_{k-1} + \log \frac{g_{\mu_1}(Y_k)}{g_{\mu_0}(Y_k)}, \, F_0 = 0. \text{ (SPRT)} \end{split}$$

- Node *I* receives *X*_{*k*,*I*} and computes *W*_{*k*,*I*}.
- Node *I* transmits *Y*_{*k*,*I*} to the Fusion node.
- Fusion node receives Y_k and computes F_k .
- Fusion node decides H_0 if $F_k \le \log \beta$, H_1 if $F_k \ge -\log \alpha$; otherwise continues.

Comments

- Noise at FC usually not considered in literature.
 - With noise, usual fusion decision rules AND, OR, Majority etc. do not work.
 - Our algorithm implicitly does it.
 - We use physical layer fusion (no MAC transmission delays)
- Noise + interference distribution at FC may not be known ⇒ Non parametric test desired.
- EMI at FC \Rightarrow Outliers will be there \Rightarrow Robustness desired.

Performance analysis of distributed algorithm

Sample Path argument



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Different algorithms at FC and secondary nodes for computing $W_{k,l}$ and F_k



- Entropy test performs best for energy detection (at secondary nodes).
- At the FC, only mean change, hence rank test/t test/random walk more appropriate.
- Rank test good only for symmetric distributions. Choose robust versions of random walk or t test.
- Conclusion: At local nodes, entropy test with energy detector.

At FC, M-random walk/M test.

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Spectrum Sensing

Gaussian noise (at FC) in the presence of Uniform outliers



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Laplacian noise (at FC) in the presence of Uniform outliers



More examples - heavy tailed without outliers Log normal noise (at FC) in the presence of Uniform outliers



More examples - heavy tailed with 1% outliers Pareto noise (at FC) in the presence of Uniform outliers



Asymptotics for distributed algorithm

N = Time when FC makes decision $P_{FA} =$ Probability of False alarm. $P_{MD} =$ Probability of Missed Detection.

$$\begin{aligned} & \Gamma_{1,l} = \rho_l |logc|, \ \Gamma_{0,l} = -\Gamma_l |logc|\\ & \beta_0 = -|logc|, \beta_1 = |logc| \end{aligned}$$

$$\begin{split} D_{tot}^{i} &= \sum_{i=1}^{L} E_{1} [X_{k,i} - \frac{\mu_{0} + \mu_{1}}{2}] \\ \zeta_{1}^{i} &= Lb_{1} + Z_{k} - \frac{\mu_{0} + \mu_{1}}{2} \\ \Delta_{i} &= Lb_{i} - \frac{\mu_{0} + \mu_{1}}{2}, \text{ i=0,1.} \\ \Gamma_{I} &= \frac{E_{0} [X_{1,I} - \frac{\mu_{0} + \mu_{1}}{2}]}{D_{tot}^{0}} \\ \rho_{I} &= \frac{E_{1} [X_{1,I} - \frac{\mu_{0} + \mu_{1}}{2}]}{D_{tot}^{1}} \end{split}$$

Theorem

- (1) $P_0(N < \infty) = 1$ and $P_1(N < \infty) = 1$ for any positive $\Gamma_{1,l}, \beta_1$ and any negative $\Gamma_{0,l}, \beta_0$.
- (2) Under H_i , $\lim_{c\to 0} \frac{N}{|logc|} \leq \frac{2}{L\lambda} + \frac{1}{\Delta_0} (1 + \frac{E_i[|\zeta_1^*|]}{D_{tot}^i})$, for i=0,1.

(3) If m.g.f. of Z_1 exists in a neighbourhood of 0, then $\lim_{c\to 0} \frac{P_{FA}}{c} = 0$ $\lim_{c\to 0} \frac{P_{MD}}{c} = 0.$

Summary

- Developed distributed, non parametric, sequential spectrum sensing algorithms for tackling SNR uncertainty issues.
- Studied the performance (via simulations and asymptotics) of different algorithms at secondary nodes for energy detection and FC MAC.
- Proposed appropriate choice of algorithms for spectrum sensing for the CR system overall.
- Extended the design to handle outliers at the FC receiver.

Thank You!