

1 Recap

- Gauss Jordan elimination

2 More on linear equations

In general, given an arbitrary row-reduced echelon matrix C , the number of non-zero rows is called the *row rank* of C . It is the number of pivots in C . It is also the dimension of the row space (HW). Bear in mind that the row space does not change under row operations.

Using the nullity-rank theorem one can prove that the row rank of C equals its *column rank* (the dimension of the column space/the number of non-zero columns in the Column reduced Column-echelon form/the number of non-zero rows in the RREF of A^T). Thus we can talk unambiguously about the *rank* of a matrix.

Returning back to $[\tilde{A}|\tilde{b}]$, the number of “free variables” equals the number of columns minus the row rank.

3 Inverses of linear maps vs matrices

An $n \times n$ (*square*) matrix A is said to be *invertible* if there exists an $n \times n$ matrix B such that $BA = AB = I$. B is called the *inverse* of A and is denoted as A^{-1} . Recall that A defines a linear map $T : \mathbb{F}^n \rightarrow \mathbb{F}^n$ as $T(v) = Av$ and the matrix of T in the standard basis of \mathbb{F}^n is A . Note that the column space of A is the range of T .

Here is an important result: The linear map T is invertible if and only if A is an invertible matrix. Moreover, the matrix associated to T^{-1} in the standard basis of \mathbb{F}^n is A^{-1} .

Proof: If T is (left-)invertible: There is a map $T^{-1} : V \rightarrow V$ such that $T^{-1}T = TT^{-1} = I$. Suppose the matrix associated to T^{-1} is B . By properties of composition, $BA = AB = I$. Hence A is invertible and B is its inverse.

If A is (left-)invertible: There is a matrix B such that $BA = I$. Hence, the corresponding linear map \tilde{T} satisfies $\tilde{T}T = I$. Since a left inverse is *the* inverse, $T\tilde{T} = I$. (As a consequence, the left inverse of the *matrix* A is its right inverse.) \square

An $n \times n$ matrix A is invertible if and only if its (column) rank is n . (Sometimes, one says “the column rank is full” or “the matrix has full rank”.) Alternatively, A is invertible if and only if $Ax = 0$ has a trivial solution.

Proof: If A is invertible: The linear map $T : \mathbb{F}^n \rightarrow \mathbb{F}^n$ defined as $T(v) = Av$ is also invertible. Hence, if e_1, \dots, e_n is a basis, then so is $T(e_1), T(e_2), \dots$. Thus, the column space is *all* of \mathbb{F}^n . Hence the rank is n .

If the rank is n : The column space is all of \mathbb{F}^n . Hence $T(e_1), T(e_2), \dots, T(e_n)$ form a basis. Therefore T is invertible. Thus, so is A .

By nullity-rank theorem, the alternative statement is true as well. \square

Later on, we shall see that the rank is full if and only if a certain expression called the ‘determinant’ is non-zero.

4 How does one compute inverses ?

The key observation is that computing an inverse is the *same* as solving a *certain* system of linear equations. This system can of course be solved (or proved to be inconsistent) using the Gauss-Jordan algorithm.

Indeed, suppose A is invertible and $[A^{-1}]_{ij} = b_{ij}$. Then $AB = I$ is equivalent to $\sum_k a_{ik}b_{kj} = \delta_{ij}$ where $\delta_{ij} = [I]_{ij}$ (the so called Kronecker delta.) In other words, for every fixed j , we have to solve a linear system for b_{1j}, b_{2j}, \dots . The other way of looking at the problem is that each *column* of B is an unknown vector x_i satisfying $Ax_i = e_i$.

Gauss-Jordan elimination to compute inverses:

- We need to form n augmented matrices $[A|e_1], [A|e_2], \dots, [A|e_n]$.
- Unless the resulting equations are *inconsistent*, that is, A is not invertible, one can bring *all* n augmented matrices to their RREFs simultaneously, by the *same* row operations. (Indeed, the A part is the same for all n matrices.)

So in practice, one applies row operations to $[A|I]$ to get $[I|A^{-1}]$. (After all, if the column rank is full, then the RREF *is* I .) Note that this procedure also lets us know whether A is invertible or not. On paper, if A is invertible, and we know the inverse A^{-1} , any linear system $Ax = b$ can be solved using $x = A^{-1}b$. However, in practice, computing the inverse is inefficient and subject to rounding-off errors.

Example: Determine if $A = \begin{bmatrix} 2 & 3 & 4 \\ 2 & 1 & 1 \\ -1 & 1 & 2 \end{bmatrix}$ is invertible. If so, find the inverse.

We must row-reduce $[A|I]$. $R_1 \leftrightarrow R_3$ and $R_1 \rightarrow -R_1$: $\left[\begin{array}{ccc|ccc} 1 & -1 & -2 & 0 & 0 & -1 \\ 2 & 1 & 1 & 0 & 1 & 0 \\ 2 & 3 & 4 & 1 & 0 & 0 \end{array} \right]$.

To clear the first column, $R_2 \rightarrow R_2 - 2R_1, R_3 \rightarrow R_3 - 2R_1$: $\left[\begin{array}{ccc|ccc} 1 & -1 & -2 & 0 & 0 & -1 \\ 0 & 3 & 5 & 0 & 1 & 2 \\ 0 & 5 & 8 & 1 & 0 & 2 \end{array} \right]$.

$R_2 \rightarrow R_2/3, R_3 \rightarrow R_3 - 5R_2, R_1 \rightarrow R_1 + R_2$: $\left[\begin{array}{ccc|ccc} 1 & 0 & -\frac{1}{3} & 0 & \frac{1}{3} & -\frac{1}{3} \\ 0 & 1 & \frac{5}{3} & 0 & \frac{1}{3} & \frac{2}{3} \\ 0 & 0 & -\frac{1}{3} & 1 & -\frac{5}{3} & -\frac{4}{3} \end{array} \right]$.

$R_3 \rightarrow -3R_3, R_2 \rightarrow R_2 - \frac{5}{3}R_3, R_1 \rightarrow R_1 + \frac{1}{3}R_3$: $\left[\begin{array}{ccc|ccc} 1 & 0 & 0 & -1 & 2 & 1 \\ 0 & 1 & 0 & 5 & -8 & -6 \\ 0 & 0 & 1 & -3 & 5 & 4 \end{array} \right]$.

Hence A is invertible and A^{-1} is $\begin{bmatrix} -1 & 2 & 1 \\ 5 & -8 & -6 \\ -3 & 5 & 4 \end{bmatrix}$.