

# Rigidity with few locations: Vertex Spanning Planar Laman Graphs in Triangulated Surfaces

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based on joint work with **Simion Tarabykin**

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# What is infinitesimal rigidity?

$G = ([n], E)$  is a finite graph. ( $[n] = \{1, 2, \dots, n\}$ .)

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$$\text{rank}(R(G, p)) = dn - \binom{d+1}{2},$$

where the *rigidity matrix*  $R(G, p) \in M_{dn \times |E|}(\mathbb{R})$  is defined by:

column  $ij \in E$  equals (say  $i < j$ )

$[0, \dots, 0, p(i) - p(j), 0, \dots, 0, p(j) - p(i), 0, \dots, 0]^T$ .

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$p$  is *generic* if its  $dn$  entries are algebraically independent over  $\mathbb{Q}$ .

Fact: for generic  $p$ ,  $\text{rank}(R(K_n, p)) = dn - \binom{d+1}{2}$ .

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$G$  is (generically) *d-rigid* if  $(G, p)$  is infinitesimally rigid for some (equivalently all) generic  $p : [n] \rightarrow \mathbb{R}^d$ .

# Examples:

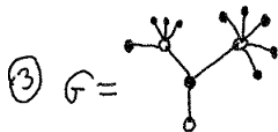
①  $G = \Delta, \mathbb{R}^2:$

$p = \text{flat triangle:}$   Non trivial infi. motion  
 $\Rightarrow (G, p)$  is not infi. rigid.

$p$  generic:  is infi. rigid.

②  $G =$    $\Rightarrow G$  is 2-rigid.

$G$  is NOT 3-rigid ( $|E| < 3 \cdot 4 - \binom{3+1}{2} = 6$ )



$$p: V \rightarrow \mathbb{R}^1$$
$$\bullet \mapsto 1$$
$$o \mapsto -1$$

$(G, p)$  is infi. rigid.

# Main objective

$G = (V, E)$  is a finite graph.

$A \subseteq \mathbb{R}^d$ .

## Definition

$G$  is *A-rigid* if there exists  $p : V \rightarrow A$  such that the framework  $(G, p)$  is infinitesimally rigid.

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$G$  is *A-rigid* if there exists  $p : V \rightarrow A$  such that the framework  $(G, p)$  is infinitesimally rigid.

$\mathcal{F}$  a family of generically  $d$ -rigid graphs.  $\mathcal{F}(n)$  is its subfamily of  $G \in \mathcal{F}$ ,  $|V(G)| \leq n$ .

## Definition

$\mathcal{F}$  is  *$d$ -rigid with  $c$  locations* if there exists  $A \subseteq \mathbb{R}^d$  of size  $c$  s.t all  $G \in \mathcal{F}$  are  $A$ -rigid.

$c_d(\mathcal{F})$  is the minimal such  $c$ .

## Main interest:

1. Families with bounded  $c_d(\mathcal{F})$ .
2. Growth of  $c_d(\mathcal{F}(n))$  as  $n \rightarrow \infty$ .



# All $d$ -rigid graphs

$\mathcal{F}_d = \{\text{all } d\text{-rigid graphs}\}.$

Fekete-Jordan 2005

$c_1(\mathcal{F}_1) = 2.$  (As a spanning tree is bipartite.)

$c_2(\mathcal{F}_2(n)) = \Omega(\sqrt{n}).$

Their argument shows: for  $d \geq 2$ ,  $c_d(\mathcal{F}_d(n)) = \Omega(\sqrt{n}).$

Sketch: let  $H$  be minimally  $d$ -rigid on  $k \geq d$  vertices.

$G = G(H)$  is obtained by: for each pair of vertices  $v, u \in H$  choose a  $d$ -subset  $B = B(v, u) \subseteq V(H)$  containing them and connect a new vertex  $v_B$  to all vertices in  $B$ .

Then  $|V(G)| = k + \binom{k}{2}$ , and in a  $d$ -rigid realization of  $G$  each vertex of  $H$  must have a different location!

Király 2021

$c_2(\mathcal{F}_2(n)) = \Theta(\sqrt{n}).$

# Planar graphs

Király 2021 (answers Whiteley)

$c_2(\text{Planar Laman}) \leq 26$ . (Laman:= edge-minimal 2-rigid)

Adiprasito-N. 2020

$c_3(\text{Maximal Planar}) \leq 76$ .

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The main algebraic statement, allowing inductive proofs in both results, is about moving vertices into  $A$ :

Adiprasito-N. 2020, also Király 2021

Assume  $(G, p)$  is infinitesimally rigid in  $\mathbb{R}^d$ ,  $\deg_G(v) = c$ ,  $A \subseteq \mathbb{R}^d$  with generic coordinates,  $|A| \geq \binom{d+c}{d}$ .

Then there exists  $a \in A$  s.t.  $(G, p')$  is infinitesimally rigid, where  $p' : V \rightarrow \mathbb{R}^d$  is defined by  $p'(v) = a$  and  $p'(u) = p(u)$  for all  $u \in V - v$ .

# Graphs on surfaces

Let  $M_g$  denotes the surface of genus  $g$  (orientable or not).

Let  $\mathcal{F}(M_g)$  be the family of graphs of triangulations of  $M_g$ .

Fogelsanger 1988

$c_3(\mathcal{F}(M_g)) \leq \aleph_0$  for all  $g$ ,

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$c_3(\mathcal{F}(M_g)) \leq c(g)$  for some constant  $c(g)$  depending on  $g$ .

Let  $\mathcal{L}(M_g)$  be the Laman graphs embedable in  $M_g$ .

Király 2021

$c_2(\mathcal{L}(M_g)) \leq c(g)$  for some constant  $c(g) = O(\sqrt{g})$ .

**Open:** are there absolute constants  $c_3, c_2$  s.t. for all  $g$ :

$$c_3(\mathcal{F}(M_g)) \leq c_3?$$

$$c_2(\mathcal{L}(M_g)) \leq c_2?$$

## Intermediate problem

Does every graph of a triangulation of a surface  $M_g$  contain a vertex spanning planar Laman subgraph?

Note: if YES then  $c_2(\mathcal{F}(M_g)) \leq 26$ .

## Theorem (N.- Simion Tarabykin)

YES if the Euler characteristic  $\chi(M_g) \geq 0$ .

## N.-Tarabykin

- Every triangulation of the projective plane  $\mathbb{R}P^2$  contains a **spanning disc**.
- Every triangulation of the Torus  $T$  contains a **spanning cylinder**.
- Every triangulation of the Klein bottle  $K$  contains a vertex spanning, planar, 2-dimensional complex; it is either a **cylinder, or a pinched disc, or a connected sum of two triangulated discs along a triangle**.

Then these vertex spanning subcomplexes are indeed 2-rigid, hence contain a spanning planar Laman subgraph, and Király's result apply.



# Irreducible triangulations

A triangulation  $\Delta$  of  $M_g$  is **irreducible** if each contraction of an edge of  $\Delta$  changes the topology; equivalently, each edge belongs to an *empty* triangle of  $\Delta$ .

Barnette-Edelson 1988/9

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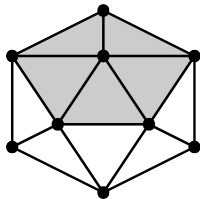
For all  $g$ ,  $M_g$  has finitely many minimal triangulations.

When  $\chi(M_g) \geq 0$  the minimal triangulations are characterized:

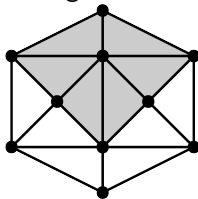
2 such  $\mathbb{R}P^2$ : Barnette 1982.

21 such  $T$ : Lavrenchenko 1990.

29 such  $K$ : Lavrenchenko-Negami 1997, Sulanke 2006.

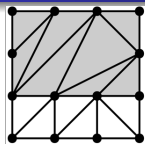


$\mathbb{R}P^2_1: (5,5,5,5,5,5)$

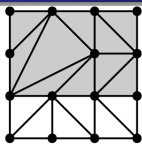


$\mathbb{R}P^2_2: (6,6,6,6,6,6,4,4,4,4)$

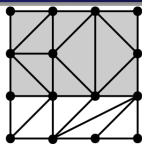
# Some spanning subcomplexes in minimal triangulations



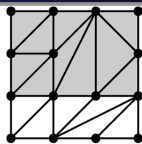
$T_1:(6,6,6,6,6,6,6)$



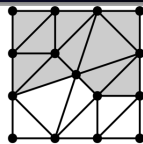
$T_2:(6,6,6,6,6,6,6,6)$



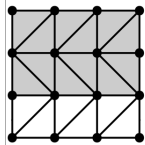
$T_3:(7,7,6,6,6,6,5,5)$



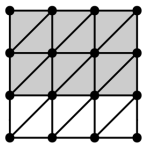
$T_4:(7,7,7,6,6,5,5,5)$



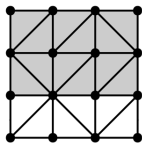
$T_5:(7,7,7,7,6,5,5,4)$



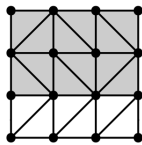
$T_6:(6,6,6,6,6,6,6,6,6,6)$



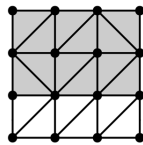
$T_7:(6,6,6,6,6,6,6,6,6,6)$



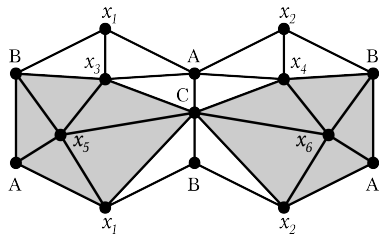
$T_8:(8,8,8,5,5,5,5,5,5,5)$



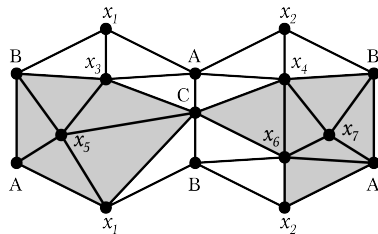
$T_9:(8,7,7,6,6,5,5,5,5,5)$



$T_{10}:(8,7,7,6,6,5,5,5,5,5)$

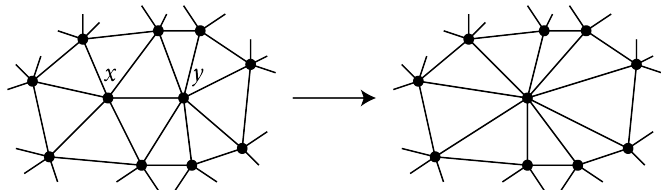


$K_{c1}:(8,8,8,5,5,5,5,5,5,5)$

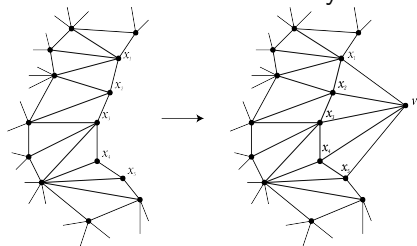


$K_{c2}:(9,9,7,6,6,5,5,5,4,4)$

# Induction via vertex splits



**Figure:** Above: edge contraction; reverse arrow for vertex split.  
Below: cone over boundary interval in the spanning subsurface.



We want the vertex split  $\Delta \rightarrow \Delta'$  to allow an **extension**  $S' \subseteq \Delta'$  of the spanning disc/cylinder/etc  $S \subseteq \Delta$ .

## Definition: extension

A spanning subsurface  $S \subseteq \Delta$  is *extendible* if for every vertex split  $\Delta \rightarrow \Delta'$  there exists  $S' \subseteq \Delta'$  s.t. either

- (i)  $S'$  is obtained from  $S$  by a split at the same vertex, or
- (ii)  $S'$  is obtained from  $S$  by coning over an interval in its boundary.

Note: then  $S' \subseteq \Delta'$  is spanning and homeomorphic to  $S$ .

## Theorem (N.-Tarabykin)

Let  $\Delta$  triangulate some  $M_g$ , and let  $S \subseteq \Delta$  be a vertex spanning subsurface. Then:

- (1)  $S$  is extendible in  $\Delta$  iff it includes at least one edge from every triangle in  $\Delta$ .
- (2) If  $S$  is extendible then it has an *extendible* extension  $S' \subseteq \Delta'$ .

# The easy direction

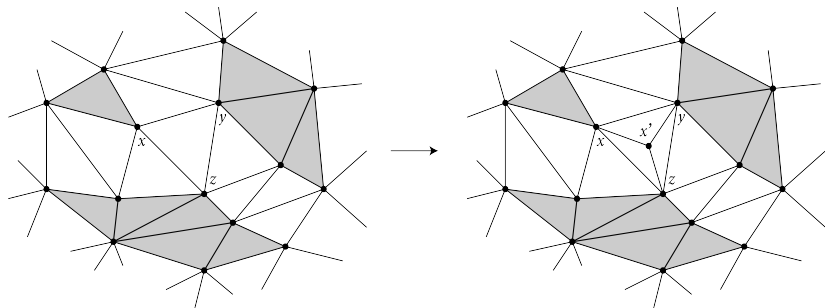


Figure: Non extendible subcomplex.

Note: all  $S$  we chose in irreducible triangulations are extendible subsurfaces, except for the 4 in the “crosscap” triangulations of the Klein bottle, which are *pinched* discs.

# How to choose an extendible $S'$ ?

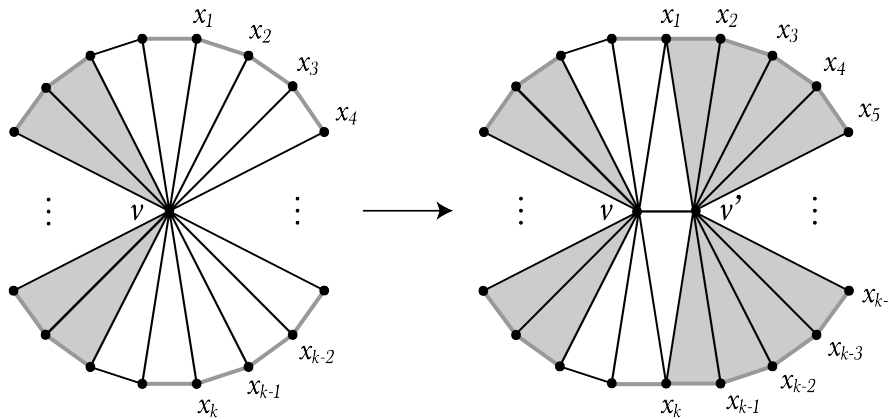


Figure:  $S \rightarrow S'$  via coning over boundary interval.

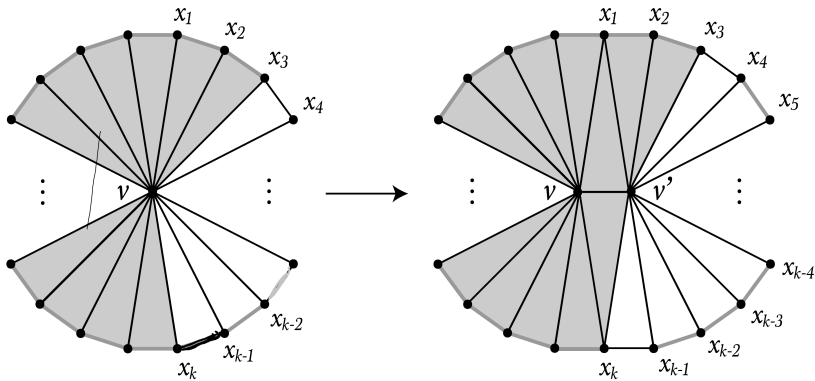
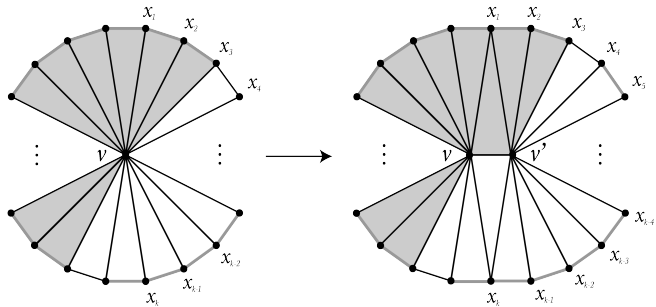
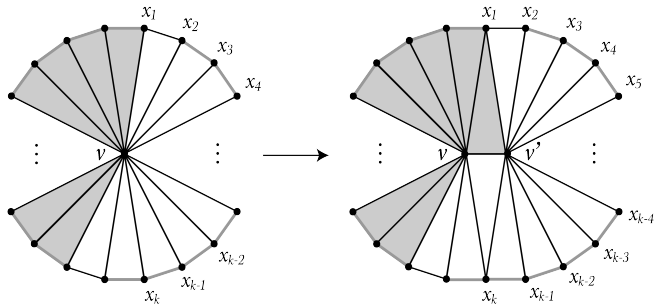


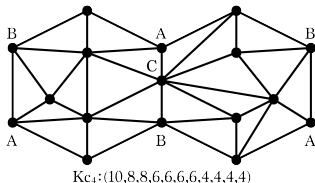
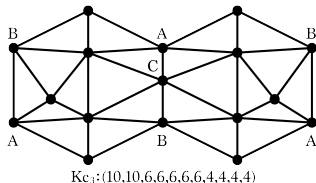
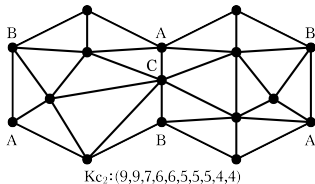
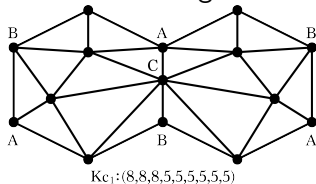
Figure:  $S \rightarrow S'$  via vertex split.





# What's left: crosscap triangulation of $K$

$ABC$  is a noncontractible cycle in each of the 4 crosscap irreducible triangulations:

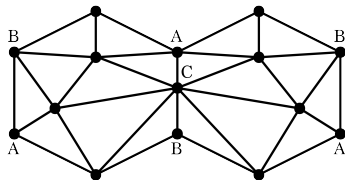
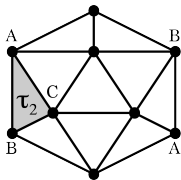
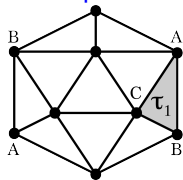


Each is a connected sum of two  $\mathbb{R}P^2$ 's along the triangle  $ABC$ ; this triangle can be part of the spanning disc in each  $\mathbb{R}P^2$ .

# Does $ABC$ survive the vertex splits?

If YES, then again the triangulation  $\Delta$  of  $K$  is a connected sum of two  $\mathbb{R}P^2$ 's along  $ABC$ , and  $ABC$  can be taken as a triangle in each spanning disc.

The connected sum of those discs is a **planar strongly connected 2-complex**.



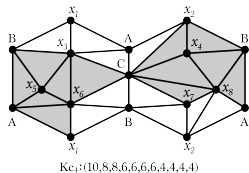
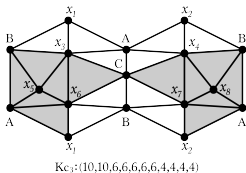
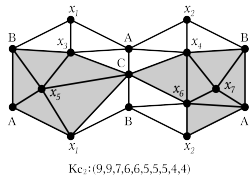
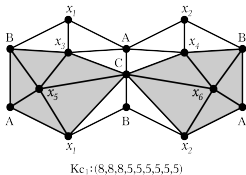
# If NOT

Then some vertex split induced a vertex split of the cycle  $ABC$ .

## Commutativity claim

The vertex splits can be rearranged s.t. the first one splits  $ABC$ .

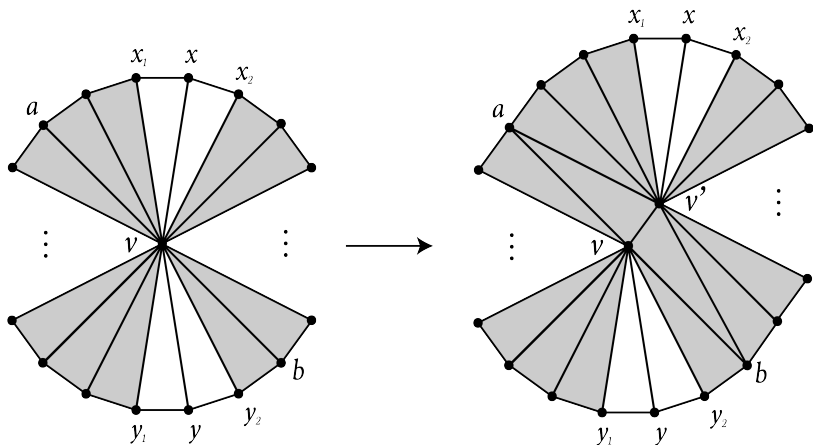
If  $C$  splits first (similary for  $A, B$ ), choose  $S$  a spanning pinched



disc at  $C$ .

# From a pinched disc to a cylinder

The first split makes  $S'$  a spanning cylinder.



**Figure:** Resolving a singularity.

The Extension Theorem shows that further splits preserve having a spanning cylinder. This completes the proof.  $\square$

# Open problems

- All surfaces: (i) Is  $c_3(\cup_g \mathcal{F}(M_g))$  finite?  
(ii) Is  $c_2(\cup_g \mathcal{L}(M_g))$  finite?  
(iii) If NO in (i), is  $c_2(\cup_g \mathcal{F}(M_g))$  finite? E.g. via:

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- What about triangulations of higher dimensional manifolds, or even just spheres?
- All  $d$ -rigid graphs,  $d > 2$ : is  $c_d(F_d(n)) = o(n)$ ?

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THANK YOU!

arXiv:2205.00558